## From sequence alignment to phylogenetic tree

Jia－Ming Chang<br>Department of Computer Science<br>National Chengchi University，Taiwan



## Nothing in biology makes sense except in the light of evolution．

－Theodosius Dobzhansky－ 1973
https：／／biologie－lernprogramme．de／daten／programme／js／homologer／daten／lit／Dobzhansky．pdf若不採用演化論，生物學的一切都說不通

## Evolution

- Charles Darwin's 1859 book (On the Origin of Species By Means of Natural Selection, or the Preservation of Favoured Races in the Struggle for Life) introduced the theory of evolution.
- At the molecular level, evolution is a process of mutation with selection.
- Molecular evolution is the study of changes in genes and proteins throughout different branches of the tree of life.
- Phylogeny is the inference of evolutionary relationships.


## 1960s: globin phylogeny

- tree of 13 orthologs by Margaret Dayhoff and colleagues
- Arrow 1: node corresponding to last common ancestor of a group of vertebrate globins.
- Arrow 2: ancestor of insect and vertebrate globins



## The neighbor-joining tree of SARS-CoV-2 related coronaviruses

- CDSs were aligned based on translated amino acid sequences using MUSCLE v3.8.31 ...
- Phylogenetic relationships were constructed using the neighbor-joining method based on Kimura's two-parameter model.



## Phylogenetic Relationship of CoVs

- Sequence alignment was carried out using MUSCLE software.
- Gblocks was used to process the gap in the aligned sequence.
- Using MegaX, we inferred all maximum likelihood phylogenetic trees.



## Flow to build Phylogenetic tree



## Lineage B clade-specific determinants for human ACE2 usage

- Replacing all 14 contact points and the surrounding amino acids (known as the receptor-binding motif (RBM)) led to increased ACE2 entry with clade 2 and 3 RBDs
- $2 \rightarrow 1$ (version 3 ) = clade 2 residues $322-400$ + clade 1 residues $400-501$
- $3 \rightarrow 1$ (version 3 ) = clade 3 residues 322-385 + clade 1 residues 386-501



## Lineage $B$ betacoronavirus RBD alignment


b

# SARS-CoV-2 RBD alignment with clade consensus RBD 



Extended Data Fig. 4 | Lineage B panel RBD sequence features. a, Amino acid sequences corresponding to SARS-spike residues 317 through 500 were aligned with ClustalW. Contact points between SARS-spike and human ACE2 are indicated with an (*). Clade 2 sequences are shown as compared to clade 2 As6526, with identical residues indicated with a (.) and sites that vary between clade 2 viruses highlighted in purple. Loop deletions are highlighted in orange. $\mathbf{b}$, Amino acid alignment of 2019-nCoV RBD and consensus RBD sequences for clade 1 and 2 and BM48-31 (clade 3). Loop deletions are highlighted in orange.

[^0]1. Sequence Alignment


By F. H. C. CRICK
Medical Research Council Unit for the Study of Molecular Biology,
Cavendish Laboratory, Cambridge

- Biologists should realise that before long we shall have a subject which might be called 'protein taxonomy'-the study of the amino acid sequences of the proteins of an organism and the comparison of them between species.
- It can be argued that these sequences are the most delicate expression possible of the phenotype of an organism and that vast amounts of evolutionary information may be hidden away within them.


## Sequence alignment

## http://phylo.cs.mcgill.ca/



### 1.1 Substitution Matrix 1. Sequence Alignment

## How Can We Compare Sequences?

- To compare Sequences, we need to compare residues
- We need to know how much it COSTS to SUBSTITUTE
- an Alanine into an Isoleucine
- a Tryptophan into a Glycine
- The table that contains the costs for all the possible substitutions is called the SUBSTITUTION MATRIX


## Making a Substitution Matrix



The Diagonal Indicates How Conserved a residue tends to be.
W is VERY Conserved

Some Residues are Easier To mutate into other similar.

## How to derive that matrix? PAM

|  | $\begin{gathered} \hline \mathrm{A} \\ \text { Ala } \\ \hline \end{gathered}$ | $\begin{array}{\|c\|} \hline \mathrm{R} \\ \mathrm{Arg} \end{array}$ | $\begin{gathered} \mathrm{N} \\ \text { Asn } \end{gathered}$ | $\begin{gathered} \mathrm{D} \\ \text { Asp } \end{gathered}$ | $\begin{gathered} \mathrm{C} \\ \mathrm{Cys} \end{gathered}$ | $\begin{gathered} \mathrm{Q} \\ \mathrm{Gln} \end{gathered}$ | $\begin{gathered} \text { E } \\ \text { Glu } \end{gathered}$ | $\begin{gathered} \mathrm{G} \\ \mathrm{Gly} \end{gathered}$ | $\begin{gathered} \hline \mathrm{H} \\ \mathrm{His} \end{gathered}$ | $\begin{gathered} \mathrm{I} \\ \mathrm{Ile} \end{gathered}$ | $\begin{gathered} \hline \text { L } \\ \text { Leu } \end{gathered}$ | $\begin{gathered} \mathrm{K} \\ \text { Lys } \end{gathered}$ | $\begin{gathered} \hline \mathbf{M} \\ \text { Met } \end{gathered}$ | $\begin{gathered} \text { F } \\ \text { Phe } \end{gathered}$ | $\begin{gathered} \hline \text { P } \\ \text { Pro } \end{gathered}$ | $\begin{gathered} \mathrm{S} \\ \mathrm{Ser} \end{gathered}$ | $\begin{gathered} \hline \text { T } \\ \text { Thr } \end{gathered}$ | $\begin{gathered} \hline \mathrm{W} \\ \mathrm{Trp} \end{gathered}$ | $\begin{gathered} \hline \mathrm{Y} \\ \mathrm{Tyr} \end{gathered}$ | $\begin{gathered} \mathrm{V} \\ \mathrm{Val} \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| R | 30 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| N | 109 | 17 |  |  |  |  |  |  |  |  |  | $f$ | mi |  | cid | - |  |  |  |  |
| D | 154 | - | 532 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| C | ${ }^{33}$ | 10 | ${ }^{\circ}$ | ${ }^{\circ}$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| Q | 93 | 120 | 50 | 76 | 0 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| E | 265 | 0 | 94 | 831 | $\bigcirc$ | 422 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| G | 579 | 10 | 156 | 162 | 10 | 30 | 112 |  |  |  |  |  |  |  |  |  |  |  |  |  |
| H | 21 | 103 | 226 | ${ }^{43}$ | 10 | 243 | ${ }^{23}$ | 10 |  |  |  |  |  |  |  |  |  |  |  |  |
| I | 66 | 30 | ${ }^{36}$ | 13 | 17 | 8 | 35 | 0 | ${ }^{3}$ |  |  |  |  |  |  |  |  |  |  |  |
| L | 95 | 17 | 37 | 0 | Y | 75 | 15 | 17 | 40 | 253 |  |  |  |  |  |  |  |  |  |  |
| K | 57 | 477 | ${ }^{322}$ | 85 | 0 | 147 | 104 | 60 | ${ }^{23}$ | 43 | 39 |  |  |  |  |  |  |  |  |  |
| M | 29 | 17 | 0 | 0 | 0 | 20 | 7 | 7 | 0 | 57 | 207 | 90 |  |  |  |  |  |  |  |  |
| F | 20 | 7 | 7 | 0 | 0 | 0 | 0 | 17 | 20 | 90 | 167 | 0 | 17 |  |  |  |  |  |  |  |
| P | 345 | 67 | 27 | 10 | 10 | 93 | 40 | 49 | 50 | 7 | 43 | 43 | 4 | 7 |  |  |  |  |  |  |
| S | 772 | 137 | 432 | 98 | 117 | 47 | 86 | 450 | 26 | 20 | 32 | 168 | 20 | 40 | 269 |  |  |  |  |  |
| T | 590 | 20 | 169 | 57 | 10 | 37 | 31 | 50 | 14 | 129 | 52 | 200 | 28 | 10 | ${ }^{73}$ | 696 |  |  |  |  |
| W | $\bigcirc$ | 27 | 3 | $\bigcirc$ | 0 | - | $\bigcirc$ | $\bigcirc$ | ${ }^{3}$ | 0 | 13 | 0 | 0 | 10 | $\bigcirc$ | 17 | $\bigcirc$ |  |  |  |
| Y | 20 | 3 | ${ }^{36}$ | 0 | ${ }^{30}$ | 0 | 10 | 0 | 40 | ${ }^{13}$ | ${ }^{23}$ | 10 | 0 | 260 | 0 | 22 | ${ }^{23}$ | 6 |  |  |
| V | 365 | 20 | 13 | 17 | 33 | 27 | 37 | 97 | 30 | 661 | 303 | 17 | 77 | 10 | 50 | 43 | 186 | 0 | 17 |  |
|  | $\begin{array}{\|l\|} \hline \text { A } \\ \text { Ala } \\ \hline \end{array}$ | $\begin{array}{\|l\|} \hline \mathrm{R} \\ \mathrm{Arg} \\ \hline \end{array}$ | $\begin{array}{\|l\|} \hline \mathrm{N} \\ \text { Asn } \\ \hline \end{array}$ | $\begin{array}{\|l\|} \hline \text { D } \\ \text { Asp } \\ \hline \end{array}$ | $\begin{aligned} & \mathrm{C} \\ & \mathrm{Cys} \\ & \hline \end{aligned}$ | $\begin{aligned} & \mathrm{Q} \\ & \mathrm{G} 1 \mathrm{n} \end{aligned}$ | $\begin{aligned} & \hline \mathrm{E} \\ & \text { Glu } \end{aligned}$ | $\begin{aligned} & \hline \text { G } \\ & \text { Gly } \end{aligned}$ | $\begin{aligned} & \hline \mathrm{H} \\ & \mathrm{His} \end{aligned}$ | $\begin{aligned} & \hline \text { I } \\ & \text { Ile } \\ & \hline \end{aligned}$ | $\begin{aligned} & \mathrm{L} \\ & \text { Leu } \end{aligned}$ | $\begin{aligned} & \hline \text { K } \\ & \text { Lys } \end{aligned}$ | $\begin{aligned} & \hline \mathrm{M} \\ & \mathrm{Met} \end{aligned}$ | $\begin{aligned} & \hline \text { F } \\ & \text { Phe } \end{aligned}$ | $\begin{aligned} & \hline \text { P } \\ & \text { Pro } \end{aligned}$ | $\begin{aligned} & \hline \text { S } \\ & \text { Ser } \end{aligned}$ | $\begin{aligned} & \hline \text { T } \\ & \text { Thr } \end{aligned}$ | $\begin{aligned} & \hline \text { W } \\ & \text { Trp } \end{aligned}$ | $\begin{aligned} & \hline \mathrm{Y} \\ & \mathrm{Tyr} \end{aligned}$ | $\begin{aligned} & \hline \mathrm{V} \\ & \mathrm{Val} \end{aligned}$ |

FIGURE 3.8 Numbers of accepted point mutations, multiplied by 10 , in 1572 cases of amino acid substitutions from closely related protein sequences. Amino acids are presented alphabetically according to the three-letter code. Notice that some substitutions (green shaded boxes) are very commonly accepted (such as V and I or S and T ). Other amino acids, such as C and W , are rarely substituted by any other residue (orange shaded boxes).

Bioinformatics and Functional Genomics, Third Edition, Jonathan Pevsner.
© 2015 John Wiley \& Sons, Ltd. Published 2015 by John Wiley \& Sons, Ltd.
Companion Website: www.wiley.com/go/pevsnerbioinformatics

## Normalized frequencies of amino acids, $f_{i}$

TABLE 3.1 Normalized frequencies of amino acid. These values sum to 1 . If the $\mathbf{2 0}$ amino acids were equally represented in proteins, these values would all be 0.05 (i.e., $5 \%$ ); instead, amino acids vary in their frequency of occurrence.

| Gly | 0.089 | Arg | 0.041 |
| :--- | :--- | :--- | :--- |
| Ala | 0.087 | Asn | 0.040 |
| Leu | 0.085 | Phe | 0.040 |
| Lys | 0.081 | Gln | 0.038 |
| Ser | 0.070 | Ile | 0.037 |
| Val | 0.065 | His | 0.034 |
| Thr | 0.058 | Cys | 0.033 |
| Pro | 0.051 | Tyr | 0.030 |
| Glu | 0.050 | Met | 0.015 |
| Asp | 0.047 | Trp | 0.010 |

some are more common $(\mathrm{G}, \mathrm{A}, \mathrm{L}, \mathrm{K})$ and some $\operatorname{rare}(\mathrm{C}, \mathrm{Y}, \mathrm{M}, \mathrm{W})$.

## The relative mutability of amino acid $j, m_{j}$

- the \# of $j$ was observed to mutate / the overall occurrence frequency of $j$ ( $f_{i}$ )
- In a scoring system alignment of two tryptophans will be weighted more heavily than two asparagines.


## TABLE 3.2 Relative mutabilities of amino acids. The value of alanine is arbitrarily set to 100.

| Asn | 134 | His | 66 |
| :--- | :--- | :--- | :--- |
| Ser | 120 | Arg | 65 |
| Asp | 106 | Lys | 56 |
| Glu | 102 | Pro | 56 |
| Ala | 100 | Gly | 49 |
| Thr | 97 | Tyr | 41 |
| Ile | 96 | Phe | 41 |
| Met | 94 | Leu | 40 |
| Gln | 93 | Cys | 20 |
| Val | 74 | Trp | 18 |

## Mutation matrix - original amino acids (columns) and replacements (rows)

- The relative mutability of amino acid $j$



## From a mutation probability matrix to a log odds matrix

$$
s_{i j}=10 \times \log _{10}\left(\frac{M_{i j}}{f_{i}}\right), \text { where } R_{i j}=\frac{M_{i j}}{f_{i}}
$$

- $M_{\mathrm{ij}}$ : models the observed change
- $f_{i}$ : the probability of a.a. $i$ occurring in the second sequence by chance
- a log scoring matrix, why?
- doing a pairwise alignment (or a BLAST search) we know what score to assign to two aligned amino acid residues.
- Logarithms are easier to use for a scoring system => sum the scores of aligned residues rather than multiply them.


# What do the numbers mean in a log odds matrix? 

- 0: neutral
- +2: indicates that the amino acid replacement occurs 1.6 times as frequently as expected by chance
- -10: that the correspondence of two amino acids in an alignment that accurately represents homology (evolutionary descent) is one tenth as frequent as the chance alignment of these amino acids


## PAM matrices

- PAM1
- At an evolutionary interval of PAM1, one change has occurred over a length of 100 amino acids.
- Other PAM matrices are extrapolated from PAM1
- PAMx = multiplied PAM1 by itself
- PAM250 matrix: for proteins that share ~20\% identity


## Mutation Matrix vs Log－odds score matrix

－Take PAM250 as an example，from asymmetric to symmetric，why？

|  |  | Original amino acid |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | A | R | N | D | C | Q | E | G | H | I | L | K | M | F | P | S | T | W | Y | V |
|  | A | 13 | 6 | 9 | 9 | 5 | 8 | 9 | 12 | 6 | 8 | 6 | 7 | 7 | 4 | 11 | 11 | 11 | 2 | 4 | 9 |
|  | R | 3 | 17 | 4 | 3 | 2 | 5 | 3 | 2 | 6 | 3 | 2 | 9 | 4 | 1 | 4 | 4 | 3 | 7 | 2 | 2 |
|  | N | 4 | 4 | 6 | 7 | 2 | 5 | 6 | 4 | 6 | 3 | 2 | 5 | 3 | 2 | 4 | 5 | 4 | 2 | 3 | 3 |
|  | D | 5 | 4 | 8 | 11 | 1 | 7 | 10 | 5 | 6 | 3 | 2 | 5 | 3 | 1 | 4 | 5 | 5 | 1 | 2 | 3 |
|  | C | 2 | 1 | 1 | 1 | 52 | 1 | 1 | 2 | 2 | 2 | 1 | 1 | 1 | 1 | 2 | 3 | 2 | 1 | 4 | 2 |
|  | Q | 3 | 5 | 5 | 6 | 1 | 10 | 7 | 3 | 7 | 2 | 3 | 5 | 3 | 1 | 4 | 3 | 3 | 1 | 2 | 3 |
| 뮹 | E | 5 | 4 | 7 | 11 | 1 | 9 | 12 | 5 | 6 | 3 | 2 | 5 | 3 | 1 | 4 | 5 | 5 | 1 | 2 | 3 |
| $0$ | G | 12 | 5 | 10 | 10 | 4 | 7 | 9 | 27 | 5 | 5 | 4 | 6 | 5 | 3 | 8 | 11 | 9 | 2 | 3 | 7 |
| 统 | H | 2 | 5 | 5 | 4 | 2 | 7 | 4 | 2 | 15 | 2 | 2 | 3 | 2 | 2 | 3 | 3 | 2 | 2 | 3 | 2 |
| 気 | I | 3 | 2 | 2 | 2 | 2 | 2 | 2 | 2 | 2 | 10 | 6 | 2 | 6 | 5 | 2 | 3 | 4 | 1 | 3 | 9 |
| $\overrightarrow{\mathrm{U}}$ | L | 6 | 4 | 4 | 3 | 2 | 6 | 4 | 3 | 5 | 15 | 34 | 4 | 20 | 13 | 5 | 4 | 6 | 6 | 7 | 13 |
| E． | K | 6 | 18 | 10 | 8 | 2 | 10 | 8 | 5 | 8 | 5 | 4 | 24 | 9 | 2 | 6 | 8 | 8 | 4 | 3 | 5 |
| － | M | 1 | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 | 2 | 3 | 2 | 6 | 2 | 1 | 1 | 1 | 1 | 1 | 2 |
| へ్ల | F | 2 | 1 | 2 | 1 | 1 | 1 | 1 | 1 | 3 | 5 | 6 | 1 | 4 | 32 | 1 | 2 | 2 | 4 | 20 | 3 |
|  | P | 7 | 5 | 5 | 4 | 3 | 5 | 4 | 5 | 5 | 3 | 3 | 4 | 3 | 2 | 20 | 6 | 5 | 1 | 2 | 4 |
|  | S | 9 | 6 | 8 | 7 | 7 | 6 | 7 | 9 | 6 | 5 | 4 | 7 | 5 | 3 | 9 | 10 | 9 | 4 | 4 | 6 |
|  | T | 8 | 5 | 6 | 6 | 4 | 5 | 5 | 6 | 4 | 6 | 4 | 6 | 5 | 3 | 6 | 8 | 11 | 2 | 3 | 6 |
|  | W | 0 | 2 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 55 | 1 | 0 |
|  | Y | 1 | 1 | 2 | 1 | 3 | 1 | 1 | 1 | 3 | 2 | 2 | 1 | 2 | 15 | 1 | 2 | 2 | 3 | 31 | 2 |
|  | V | 7 | 4 | 4 | 4 | 4 | 4 | 4 | 5 | 4 | 15 | 10 | 4 | 10 | 5 | 5 | 5 | 7 | 2 | 4 | 17 |



FIGURE 3．14 Log－odds matrix for PAM250．High PAM values（e．g．，PAM250）are useful for align－ ing very divergent sequences．A variety of algorithms for pairwise alignment，multiple sequence align－ ment，and database searching（e．g．，BLAST）allow you to select an assortment of PAM matrices such as PAM250，PAM70，and PAM30．Adapted from NCBI，ftp：／／ftp．ncbi．nlm．nih．gov／blast／matrices／．

## Why does PAM1 become symmetric?

- $M_{i j}=\frac{\lambda m_{j} A_{i j}}{\sum_{i=1, i \neq j}^{20} A_{i j}}=m_{j} \times \frac{\lambda A_{i j}}{\sum_{i=1, i \neq j}^{20} A_{i j}}=\frac{\sum_{i=1, i \neq j}^{20} A_{i j}}{f_{j}} \times \frac{\lambda A_{i j}}{\sum_{i=1, i \neq j}^{20} A_{i j}}=\frac{\lambda A_{i j}}{f_{j}}$
- $M_{j i}=\frac{\lambda m_{i} A_{i j}}{\sum_{i=1, i \neq j}^{20} A_{i j}}=m_{i} \times \frac{\lambda A_{i j}}{\sum_{i=1, i \neq j}^{20} A_{i j}}=\frac{\sum_{i=1, i \neq j}^{20} A_{i j}}{f_{i}} \times \frac{\lambda A_{i j}}{\sum_{i=1, i \neq j}^{20} A_{i j}}=\frac{\lambda A_{i j}}{f_{i}}$
- $R_{i j}=\frac{M_{i j}}{f_{i}}=\frac{\frac{\lambda A_{i j}}{f_{j}}}{f_{i}}=\frac{\lambda A_{i j}}{f_{i} f_{j}}=\frac{\frac{\lambda A_{i j}}{f_{i}}}{f_{j}}=\frac{M_{j i}}{f_{j}}=R_{j i}$


# BLOcks SUbstitution Matrix (BLOSUM) 

Henikoff, S.; Henikoff, J.G. (1992). "Amino Acid Substitution Matrices from Protein Blocks". PNAS. 89 (22): 10915-10919

## Procedure of BLOSUM

- Cluster together sequences in a family whenever more than L\% identical residues are shared, for BLOSUM-L.
- Based on local alignments \& use aligned ungapped regions of protein families.
- Count number of substitutions across different clusters (in the same family).
- Estimate frequencies using the counts.


## Summary of PAM and BLOSUM matrices

- BLOSUM62 is a matrix calculated from comparisons of sequences with no less than 62\% divergence.
- the default matrix in BLAST 2.0
- Most widely used (PAM250)
- A higher PAM number, and a lower BLOSUM number, tends to correspond to a matrix tuned to more divergent proteins.

BLOSUM90
BLOSUM62
PAM30

Less divergent
PAM120
PAM250

Human versus
chimpanzee beta globin
1.2 Pairwise Alignment

## HOW Can we Align Two Sequences?

Different types of pairwise comparisons

Method name
|Dot-plot

Local alignments

Global alignments

## Situation

General exploration of your sequence Discovering repeats
Finding long insertion/deletions
Extracting portions of sequences to make a multiple alignment

Comparing sequences with partial homology Making high quality alignments Making residue-per-residue analysis

Comparing two sequences over their entire length Identifying long insertion/deletions Checking the quality of your data Identifying every mutation in your sequences

## Global Alignments

- Take 2 Nice Protein Sequences
- A good Substitution Matrix (Blosum62)
- DYNAMIC PROGRAMMING
$>$ Seq1
THEFATCAT
$>$ Seq 2
THEFASTCAT



## Dynamic Programming

## THE THEORY OF DYNAMIC PROGRAMMING

## RICHARD BELLMAN

1. Introduction. Before turning to a discussion of some representative problems which will permit us to exhibit various mathematical features of the theory, let us present a brief survey of the fundamental concepts, hopes, and aspirations of dynamic programming.

To begin with, the theory was created to treat the mathematical problems arising from the study of various multi-stage decision processes, which may roughly be described in the following way: We have a physical system whose state at any time $t$ is determined by a set of quantities which we call state parameters, or state variables. At certain times, which may be prescribed in advance, or which may be determined by the process itself, we are called upon to make decisions which will affect the state of the system. These decisions are equivalent to transformations of the state variables, the choice of a decision being identical with the choice of a transformation. The outcome of the preceding decisions is to be used to guide the shoice of future ones, with the purpose of the whole process that of maximizing some function of the parameters describing the final state.

## Using Dynamic Programming To Align Sequences

- DP invented in the 1950s by Bellman
- Programming $\Leftrightarrow$ Tabulation
- Re-invented in 1970 by Needlman and Wunsch
- It took 10 year to find out...
- Needleman, Saul B. \& Wunsch, Christian D. (1970). "A general method applicable to the search for similarities in the amino acid sequence of two proteins". Journal of Molecular Biology. 48 (3): 443-53


## Global Alignment

Needleman, Saul B. \& Wunsch, Christian D. (1970). "A general method applicable to the search for similarities in the amino acid sequence of two proteins". Journal of Molecular Biology. 48 (3): 443-53

## The Principal of DP

- If you extend optimally an optimal alignment of two sub-sequences, the result remains an optimal alignment



## Finding the score of $i, j$

- Sequence 1: [1-i]
- Sequence 2: [1-j]
- The optimal alignment of [1-i] vs [1-j] can finish in three different manners:



## Finding the score of $i, j$



## Formalizing the algorithm



## Arranging Everything in a Table

|  | - | $F$ | $A$ | $T$ |
| :---: | :---: | :---: | :---: | :---: |
| - |  |  |  |  |
| $F$ |  | $1 \ldots . \mathrm{I}-1$ <br> $1 \ldots-1$ | $1 \ldots \mathrm{I}$ <br> $1 \ldots \mathrm{~J}-1$ |  |
| $A$ |  | $1 \ldots . \mathrm{I}-1$ <br> $1 \ldots \mathrm{~J}$ | $1 \ldots \mathrm{I}$ <br> $1 \ldots \mathrm{~J}$ |  |
| S |  |  |  |  |
| T |  |  |  |  |

## Filing Up The Matrix



|  | - | $F$ | $A$ | $T$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | 0 |  | -1 | -2 | -3 |
| $F$ | -1 | +2 | -2 | -2 | -3 |

## Delivering the alignment: Traceback

|  | - | F | A | T |
| :---: | :---: | :---: | :---: | :---: |
| - | 0 |  |  |  |
| F |  |  |  |  |
| A |  |  |  |  |
| S | $-3$ |  |  |  |
| T | -4 |  |  |  |


| T | S | A | F |
| :--- | :--- | :--- | :--- |
| T | - | A | F |

Score of 1... 3 Vs 1... 4
$\Leftrightarrow$
Optimal Aln Score

## Trace-back: possible implementation

```
while (!(i==0 && j==0)):
    if (direc_m[i][j]== 'sub'): #SUBSTITUTION
            aln1[aln_len]=pro1Seq[--i]
    aln2[aln_len]=pro2Seq[--j]
    elif (direc_m[i][j]=='del'): #DELETION
        aln1[aln_len]=' -'
        aln2[aln_len]=pro2Seq[--j]
    elif (direc_m[i][j]=='ins'): #INSERTION
    aln1[aln_len]=pro1Seq[0][--i]
    aln2[aln_len]='-'
    aln_len++
}
```


## Local Alignment Smith \& Waterman algorithm

Smith, T. F. \& Waterman, M. S. Identification of common molecular subsequences. J. Mol. Biol. 147, 195-7 (1981).

## Global alignment VS local alignment

- Global : extends from one end of each sequence to the other.
- Local : finds optimally matching regions within two sequences,
- Subsequences useful to find domains (or limited regions of homology) within sequences
- Smith and Waterman (1981) solved the problem of performing optimal local sequence alignment.
- Other methods (BLAST, FASTA) are faster but less thorough.


GLOBAL Alignment


LOCAL Alignment

## Global alignment (top) includes matches ignored by local alignment (bottom)

(a)

| NP_824492.1 | 1 | MCGDMTVHTVEYIRYR IPEQQSAEFLAAYTRAAAQLAAAPQCVDYELARC |
| :---: | :---: | :---: |
| NP_337032.1 | 1 |  |
| NP_824492.1 | 51 | EEDFEHFVLRITWTSTEDHIEGFRKSELFPDFLAEIRPYISSIEEMRHYK |
| NP_337032.1 | 1 |  |
| NP_824492.1 | 101 | PTTVRGTGAAVPTLYAWAGGAEAFARLTEVFYBKVLKDDVLAPVFEGMAP <br> $:$ \| . . . . . . | . . |||:. | . : . . .||. . | .: |:|| . |: |
| NP_337032.1 | 1 | MEGMDQMPKS FYDAVGGAKTFDAIVSRFYAQVAEDEVLRRVY----P |
| NP_824492.1 | 151 |  |
| NP_337032.1 | 44 | EDDLAGAEERLRMFLEQYWGGPRTYSE-QRGHPRLRMRHAPFRISLIERD |
| NP_824492.1 | 196 | RWVNLLQDAADDAGLPT-DAEFRSAFLAYAEWGTRLAVYFSGPDAVPPAE |
| NP_337032.1 | 93 | AVLRCMHTAVASIDSETLDDEHRRELLDYLEMAAHSLV--NSPF |
| NP_824492.1 | 245 | QPVPQWSWGAMPPYQP 260 |
| NP_337032.1 | 135 | 134 |

(b)

NP_824492.1

NP_337032.1
NP_824492.1
NP_337032.1
NP_824492.1
113 TLYAWAGGABAFARLTEVFYEKVLKDDVLAPVFEGMAPEH --..- AAHVA

SFYDAVGGAKTFDAIVSRFYAOVAEDEVLRRVY----PEDDLAGAEERLR
158 LWLGEVFGGPAAYSETQGGHGHMVAKHLGKNITEVQRRRWVNLLQDAADD

NP_337032.1
208 AGLPT-DAEFRSAFLAYAE 225

Local: $30 \%$ identity

## Global: <br> 15\% identity

## The Smith and Waterman Algorithm

- $0=>$ Ignore the rest of the Matrix $=>$ terminate a local alignment



## Filing Up a SW Matrix

$F(\mathrm{i}, \mathrm{j})=$ best

$$
\left\{\begin{array}{l}
F(\mathrm{i}-1, \mathrm{j})+\text { Gep } \begin{array}{l}
\begin{array}{l}
1 \ldots . . \mathrm{i} \\
1 \ldots . \mathrm{j}-1
\end{array}
\end{array}+\begin{array}{|}
- \\
\mathrm{x}
\end{array} \\
F(\mathrm{i}-1, \mathrm{j}-1)+\text { Mat[i,j]} \begin{array}{l}
1 \ldots . . \mathrm{i}-1 \\
1 \ldots \mathrm{j}-1
\end{array}
\end{array}+\begin{array}{l}
\mathrm{x} \\
\mathrm{x}
\end{array}\right]
$$



## Filling up a SW matrix: borders



## Filling up a SW matrix

- Best Local score $\Leftrightarrow$ Beginning of the trace-back

| * | - | A | N | I | C | E | C | A | T |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| - | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| C | 0 | 0 | 0 | 0 | 2 | 0 | 2 | 0 | 0 |
| A | 0 | 2 | 0 | 0 | 0 | 0 | 0 | 4 | 0 |
| T | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 2 | 6 |
| A | 0 | 2 | 0 | 0 | 0 | 0 | 0 | 0 | 4 |
| N | 0 | 0 | 4 | 2 | 0 | 0 | 0 | 0 | 2 |
| D | 0 | 0 | 2 | 2 | 0 | 0 | 0 | 0 | 0 |
| O | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| G | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## Adding Affine Gap Penalties <br> Forcing a bit of Biology into your alignment

Gotoh, O. An improved algorithm for matching biological sequences. J. Mol. Biol. 162, 705-8 (1982).

## Gap Penalties: Opening \& extension

- Gaps : Positions at which a letter is paired with a null are called.
- Gap scores are typically negative.
- Opening a gap is more expensive than extending it
- Since a single mutational event may cause the insertion or deletion of more than one residue, the presence of a gap is ascribed more significance than the length of the gap.
- Thus there are separate penalties for gap open and gap extension.

> Gap Opening Penalty
> Gap Extension Penalty

Seq A GARFIELDTHE----CAT
$|||||||||\mid$
Seq BGARFIELDTHELASTCAT

## But Harder To compute …

- More Than 3 Ways to extend an Alignment



## More Questions Need to be asked

- For instance, what is the cost of an insertion ?



## Solution: Maintain 3 Tables

$M(i, j)=$ best $\quad\left\{\left.\begin{array}{l}M(i-1, j-1)+\operatorname{Mat}(i, j) \\ I x(i-1, j-1)+\operatorname{Mat}(i, j) \\ \operatorname{Iy}(i-1, j-1)+\operatorname{Mat}(i, j)\end{array} \begin{array}{|l|}1 \ldots . . i-1 \\ 1 \ldots . j-1\end{array} \right\rvert\, \begin{array}{l}X \\ X\end{array}\right.$


## A Score in Linear Space

- You never Need More Than The Previous Row To Compute the optimal score



## A Score in Linear Space

for $\mathrm{i}=1$ :
for $\mathrm{j}=1$ : J
R2[i][j]=best
R2[j-1], +gep
R1[j-1]+mat R1[j]+gep
for J,
R1[j]=R2[j]

## A Score in Linear Space



You never Need More Than The Previous Row To Compute the optimal score
You only need the matrix for the TraceBack,

## Or do you ????

## An Alignment in Linear Space

Forward Algorithm

$B(i, j)+F(i, j)=$
Optimal score of the alignment that passes through pair $\mathrm{i}, \mathrm{j}$

Backward algorithm

## An Alignment in Linear Space



Backward algorithm

Forward Algorithm


Optimal B(i,j)+F(i,j)

Backward algorithm

## An Alignment in Linear Space

Forward Algorithm


Backward algorithm


Recursive divide and conquer strategy: Myers and Miller (Durbin p35)

## Remember Not To Run Out of Memory

－A survey paper
－Chao，K．－M．，Hardison R．C．and Miller，W．，1994，Recent Developments in Linear－Space Alignment Methods：a Survey，Journal of Computational Biology，1：271－291．


趙坤茂（Kun－Mao Chao）台大資工系

## Recap: Pairwise alignment

- Needleman and Wunsch: Delivers the best scoring global alignment
- Smith and Waterman: NW with an extra state 0

- Affine Gap Penalties: Making DP more realistic

- Linear space: Using Divide and Conquer Strategies Not to run out of memory
1.3 Multiple Sequence Alignment


## Sometimes two sequences are not enough...

- The man with TWO watches NEVER knows the time



## The COMPUTATIONAL Problem

- A nice set of Sequences
- Substitution Matrix (Blosum)
- Gap Penalties
- An Evaluation/Scoring Function
- An Alignment Algorithm


## What is A Multiple Sequence Alignment?

- Structural Criteria
- Residues are arranged so that those playing a similar role end up in the same column.
- Evolution Criteria
- Residues are arranged so that those having the same ancestor end up in the same column.

```
chite ---ADKPKRPLSAYMLWLNSARESIKRENPDFK-VTEVAKKGGELWRGLKD
wheat --DPNKPKRAPSAFFVFMGEFREEFKQKNPKNKSVAAVGKAAGERWKSLSE
trybr KKDSNAPKRAMTSFMFFSSDFRS----KHSDLS-IVEMSKAAGAAWKELGP
mouse -----KPKRPRSAYNIYVSESFQ----EAKDDS-AQGKLKLVNEAWKNLSP
```

chite AATAKQNYIRALQEYERNGG-
wheat ANKLKGEYNKAIAAYNKGESA
trybr AEKDKERYKREM---------
mouse AKDDRIRYDNEMKSWEEQMAE

* : .* . :


By peellden－自己的作品
By Rico Heil（User：Silmaril）－private photo

## Scoring function

- Sum of Pair (SP)
- Tree Cost: MSA with tree cost will be called tree alignment.
- Circular Sum(CS)
$\left.\left.\begin{array}{ll}\mathbf{S}_{1}: \mathbf{A T T C G} \\ \mathbf{S}_{\mathbf{2}}: \mathbf{A G T C G} \\ \mathbf{S}_{\mathbf{3}}: \mathbf{A T C A G} & \mathrm{MSA}\end{array} \begin{array}{l}\mathbf{S}_{\mathbf{1}}^{\prime}: \mathbf{A T}-\mathbf{T C}-\mathbf{G} \\ \mathbf{S}_{\mathbf{2}}^{\prime}: \mathbf{A}-\mathbf{G T C}-\mathbf{G} \\ \mathbf{S}_{\mathbf{3}}^{\prime}: \mathbf{A T}--\mathbf{C A G}\end{array}\right\} \begin{array}{l}2 \\ 4\end{array}\right\} 2$

Cost $=8$

## MSA with SP-Score: Exact Algorithm

- Given
- $k$ : \# of Sequences
- $n$ : Sequences of length
- Exactly by Dynamic Programming
- $O\left(2 n^{k}\right)$ : D.Snakoff, Simultaneous solution of RNA folding, alignment and Protosequence prolblems, SIAM J. Appl. Math.,(1985)
- Exact methods of multiple alignment use dynamic programming and are guaranteed to find optimal solutions. But they are not feasible for more than a few sequences.


## MSA with SP-Score: Complexity

- Wang L. Jiang T. On the complexity of multiple sequence alignment, J Comput Biol 1994 Winter;1(4):337-48
- multiple alignment with SP-Score => NP-complete reduction from shorest common supersequence (non-metric : not symmetry)

Table 1. Score Scheme I

| $S$ | 0 | 1 | $a$ | $b$ | $\Delta$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 2 | 2 | 1 | 2 | 1 |
| 1 | 2 | 2 | 2 | 1 | 1 |
| $a$ | 1 | 2 | 0 | 2 | 1 |
| $b$ | 2 | 1 | 2 | 0 | 1 |
| $\Delta$ | 1 | 1 | 1 | 1 | 0 |

- multiple tree alignment $=>$ MAX SNP-hard
- Paola Bonizzoni, Gianluca Della Vedoa The complexity with Multiple sequence alignment with SP-score that is a metric, Theoretical Computer Science; 259 (2001) 63-79
- multiple alignment with SP-Score => NP-complete reduction from node cover


# Feng-Doolittle algorithm 

D.F.Feng, R.F.Doolittle, Progressive sequence alignment as a prerequisite to correct phylogenetic trees. J. Mol. Evol. 25, 351-360., (1987)

## Progressive alignment

- Any exact method would be TOO SLOW.
- We will use a heuristic algorithm.
- Progressive alignment algorithm is the most popular
- use a guide tree (related to a phylogenetic tree) to determine how to combine pairwise alignments one by one to create a multiple alignment.
- Examples
- ClustalW
- MUSCLE
-     - Greedy Heuristic (No Guaranty)
-     + Fast


## Feng-Doolittle MSA occurs in 3 stages

- Feng and Dolittle, 1988; Taylor 1989

1. Do a set of global pairwise alignments

- Needleman and Wunsch' s dynamic programming algorithm

2. Create a guide tree
3. Progressively align the sequences


## Generate global pairwise alignments (Progressive 1/3)

| SeqA | Name | Len(aa) | SeqB | Name | Len(aa) | Score |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $=======================================================$ |  |  |  |  |  |  |
| 1 | beta_globin | 147 | 2 | mYoglobin | 154 | 25 |
| 1 | beta_globin | 147 | 3 | neuroglobin | 151 | 15 |
| 1 | beta_globin | 147 | 4 | soybean | 144 | 13 |
| 1 | beta_globin | 147 | 5 | rice | 166 | 21 |
| 2 | myoglobin | 154 | 3 | neuroglobin | 151 | 16 |
| 2 | myoglobin | 154 | 4 | soybean | 144 | 8 |
| 2 | myoglobin | 154 | 5 | rice | 166 | 12 |
| 3 | neuroglobin | 151 | 4 | soybean | 144 | 17 |
| 3 | neuroglobin | 151 | 5 | rice | 166 | 18 |
| 4 | soybean | 144 | 5 | rice | 166 | 43 |
| nict score |  |  |  |  |  |  |

## Guide tree (Progressive 2/3)

- Convert similarity scores to distance scores
- A tree shows the distance between objects
- Use UPGMA (defined in the phylogeny chapter)
- ClustalW provides a syntax to describe the tree

```
!
beta_globin:0.36022,
myoglobin:0.38808,
!
neuroglobin:0.39924,
(
soybean:0.30760,
rice:0.26184)
:0.13652)
:0.06560);
```

Progressive alignment (Progressive 3/3)

- Make a MSA based on the order in the guide tree
- Start with the two most closely related sequences
- Then add the next closest sequence
- Continue until all sequences are added to the MSA
- Rule: once a gap, always a gap, why?
- Gaps are often added to the first two (closest) sequences
- To change the initial gap choices later on would be to give more weight to distantly related sequences
- To maintain the initial gap choices is to trust that those gaps are most believable


## Progressive Alignment



Dynamic Programming Using A Substitution Matrix


## ClustalW

Thompson, J. D., Higgins, D. G. \& Gibson, T. J. CLUSTAL W: improving the sensitivity of progressive multiple sequence alignment through sequence weighting, position-specific gap penalties and weight matrix choice. Nucleic Acids Res. 22, 4673-80 (1994).

## The top 100 papers



Credit by Nature, http://www.nature.com/news/the-top-100-papers-1.16224

## The top 100 papers

Click through to explore the Web of Science's all-time top-cited papers. (Data provided by Thomson Reuters, extracted on 7
October 2014).

## Rank: 10 Citations: 40,289

Clustal W: improving the sensitivity of progressive multiple sequence alignment through sequence weighting, position-specific gap penalties and weight matrix choice.
Thompson, J. D., Higgins, D. G. \& Gibson, T. J
Nucleic Acids Res. 22, 4673-4680 (1994).


Credit by The top 100 papers. Van Noorden R, Maher B, Nuzzo R. Nature. 2014 Oct 30;514(7524):550-3 sequence weighting, position-specific gap penalties and weight matrix choice JD Thompson, DG Higgins, TJ Gibson Nucleic acids research 22 (22), 4673

T-Coffee: A novel method for fast and accurate multiple sequence alignment
C Notredame, DG Higgins, J Heringa
Journal of molecular biology 302 (1), 205-217


Multiple Sequence Alignment

## ClustaIW

But owing to the vagaries of citation habits, BLAST has been bumped down the list by Clustal, a complementary programme for aligning multiple sequences at once. Clustal allows researchers to describe the evolutionary relationships between sequences from different organisms, to find matches among seemingly unrelated sequences and to predict how a change at a specific point in a gene or protein might affect its function. A 1994 paper describing ClustalW, a user-friendly version of the software, is currently number 10 on the list. A 1997 paper on a later version called ClustalX is number 28.

## ClustaIW

The team that developed ClustalW, at the European Molecular Biology Laboratory in Heidelberg, Germany, had created the program to work on a personal computer, rather than a mainframe. But the software was transformed when Julie Thompson, a computer scientist from the private sector, joined the lab in 1991. "It was a program written by biologists; I'm trying to find a nice way to say that," says Thompson, who is now at the Institute of Genetics and Molecular and Cellular Biology in Strasbourg, France. Thompson rewrote the program to ready it for the volume and complexity of the genome data being generated at the time, while also making it easier to use.

The teams behind BLAST and Clustal are competitive about the ranking of their papers. It is a friendly sort of competition, however, says Des Higgins, a biologist at University College Dublin, and a member of the Clustal team. "BLAST was a gamechanger, and they've earned every citation that they get."

## Thompson et al. (1994) for an explanation of the three stages of progressive alignment implemented in ClustalW



Figure 1. The basic progressive alignment procedure, illustrated using a set of 7 globins of known tertiary structure. The sequence names are from Swiss Prot (38): Hbe_Horse: horse $\alpha$-globin; Hbe_Human: human $\alpha$-globin; Hbb_Horse: horse $\beta$-globin; Hbb__Human: human $\beta$-globin; Myg_Phyca: sperm whale myoglobin; GIb5_Petma: lamprey cyanohaemoglobin; Lgb2_Luplu: lupin leghaemoglobin. In the distance matrix, the mean number of differences per residue is given. The unrooted tree shows all branch lengths drawn to scale. In the rooted tree, all branch lengths (mean number of differences per residue along each branch) approximate positions of the $7 \alpha$-helices common to all 7 proteins are shown. This alignment was derived using CLUSTAL W with default parameters and the PAM (3) series of weight matrices.

In Figure 1 we give the $7 \times 7$ distance matrix between the 7 globin sequences calculated using the full dynamic programming method.

The guide tree
The trees used to guide the final multiple alignment process are The trees used to guide the final multiple alignment process are
calce the distance matrix of step 1 using the Neighbourcalculated from the distance matrix of step 1 using the Neighbour-
Joining method (21). This produces unrooted trees with branch Joining method (21). This produces unrooted trees with branch
lengths proportional to estimated divergence along each branch. lengths proportional to estimated divergence along each branch.
The root is placed by a 'mid-point' method (15) at a position The root is placed by a 'mid-point' method (15) at a position where the means of the branch lengths on either side of the root are equal. These trees are also used to derive a weight for each sequence (15). The weights are depenhich have a common branch with other sequences share the weight derived from the shared branch. In the example in Figure 1, the leghaemoglobin (Lgb2_Luplu) gets a weight of 0.442 , which is equal to the length of the branch from the root to it. The human $\beta$-globin (Hbb_Human) gets a weight consisting of the length of the branch leading to it that is not shared with any other sequences 0.081 ) plus half the length of the branch shared with the horse $\beta$-globin ( $0.226 / 2$ ) plus one quarter the length of the branch shared by all four haemoglobins $(0.061 / 4)$ plus one fifth the branch shared between the haemoglobins and myoglobin ( $0.015 / 5$ ) plus one sixth the branch leading to all the vertebrate globins ( 0.062 ). This sums to a total of 0.221 . In contrast, in the normal progressive alignment algorithm, all sequences would be equally weighted. The rooted tree with branch lengths and sequence weights for the 7 globins is given in Figure 1.

## Progressive alignment

The basic procedure at this stage is to use a series of pairwise alignments to align larger and larger groups of sequences following the branching order in the guide tree. You proceed from the tips of the rooted tree towards the root. In the globin example in Figure 1 you align the sequences in the following order: human vs. horse $\beta$-globin; human vs. horse $\alpha$-globin; the $2 \alpha$-globins vs. the $2 \beta$-globins; the myoglobin vs. the haemoglobins; the cyanohaemoglobin vs. the haemoglobins plus myoglobin; the leghaemoglobin vs. all the rest. At each stage a full dynamic programming $(26,27)$ algorithm is used with residue weight matrix and penalties for opening and extending gaps. Each step consists of aligning two existing alignments or sequences. Gaps that are present in older alignments remain fixed. In the basic algorithm, new gaps that are introduced at each stage

## Pairwise alignment:

Calculate distance matrix


Unrooted neighborjoining tree





## Additional features of ClustalW improve its ability to generate accurate MSAs

- Individual weights are assigned to sequences; very closely related sequences are given less weight, while distantly related sequences are given more weight
- Scoring matrices are varied dependent on the presence of conserved or divergent sequences, e.g.:
- PAM20 80-100\% id
- PAM60 60-80\% id
- PAM120 40-60\% id
- PAM350 0-40\% id
- Residue-specific gap penalties are applied


## Iterative approaches

## Iterative methods

- compute a sub-optimal solution and keep modifying that intelligently using dynamic programming or other methods until the solution converges.
- MUSCLE, Mafft, HMMs, HMMER, SAM,, IterAlign, Praline
- +: Good Profile Generators
- -: Slow, Sometimes Inaccurate



## Muscle

- Edgar, R. MUSCLE: a multiple sequence alignment method with reduced time and space complexity. Bmc Bioinformatics 5, 1-19 (2004).
- Edgar, R. MUSCLE: multiple sequence alignment with high accuracy and high throughput. Nucleic Acids Res 32,1792-1797 (2004)



## MUSCLE: Improve the progressive alignment

- Build a draft progressive alignment
- Determine pairwise similarity through $k$-mer counting
- Compute triangular distance matrix
- Construct tree using UPGMA
- Construct draft progressive alignment following tree
- Compute pairwise identity through current MSA

- Construct new tree with Kimura distance measures
- Compare new and old trees: if improved, repeat this step, if not improved, then we're done
- Refinement of the MSA
- Split tree in half by deleting one edge
- Make profiles of each half of the tree
- Re-align the profiles
- Accept/reject the new alignment


## MAFFT : Fast Fourrier Transforme

- Katoh, K., Misawa, K., Kuma, K. \& Miyata, T. MAFFT: a novel method for rapid multiple sequence alignment based on fast Fourier transform. Nucleic Acids Res. 30, 3059-66 (2002).
- Katoh, K. \& Toh, H. Recent developments in the MAFFT multiple sequence alignment program. Brief. Bioinformatics 9, 286-98 (2008).



## MAFFT

- Uses Fast Fourier Transform to speed up profile alignment
- Uses fast two-stage method for building alignments using $k$-mer frequencies
- Offers many different scoring and aligning techniques
- One of the more accurate programs available
- Available as standalone or web interface
- Many output formats, including interactive phylogenetic trees


## Iterative method of MAFFT



## Consistency-based approaches

## Progressive Alignment When It Doesn' t Work


CLUSTALW (Score=20, Gop=-1, Gep=0, $\mathrm{M}=1$ )
SeqA GARFIELD THE LAST FA-T CAT
SeqB GARFIELD THE FAST CA-T ---
SeqC GARFIELD THE VERY FAST CAT
SeqD -------- THE ---- FA-T CAT

SeqC GARFIELD THE VERY FAST CAT
CORRECT (Score=24)
SeqA GARFIELD THE LAST FA-T CAT
SeqB GARFIELD THE FAST ---- CAT
SeqC GARFIELD THE VERY FAST CAT
SeqD -------- THE ---- FA-T CAT

Multiple sequence alignment: consistency

- generally use a database of both local highscoring alignments and long-range global alignments to create a final alignment
- These are very powerful and very accurate methods
- Examples: T-Coffee, Prrp, DiAlign, ProbCons


## Mixing Heterogenous Data With T-Coffee



## www.tcoffee.org

## TCOFFEE

## T-Coffee

A collection of tools for Computing, Evaluating and Manipulating Multiple Alignments of DNA, RNA, Protein Sequences and Structures

Alignment

T-Coffee Aligns DNA, RNA or Proteins using the default T-Coffee >> Cite
M-Coffee Aligns DNA, RNA or Proteins by combining the output of popular aligners >> Cite
R-Coffee Aligns RNA sequences using predicted secondary structures $\gg$ Cite
Expresso Aligns protein sequences using structural information $\gg$ Cite
PSI-Coffee Aligns distantly related proteins using homology extension (slow and accurate) >> Dite
TM-Coffee Aligns transmembrane proteins using homology extension 패페 > Cite
Pro-Coffee Aligns homologous promoter regions Mem Clite
Accurate Automatically combine the most accurate modes for DNA, RNA and Proteins (experimental!)
Combine Combines two (or more) multiple sequence alignments into a single one $\gg$ Cite

Evaluation

Core Evaluates your Alignment and outputs a Colored version indicating the local reliability. >> Cite iRMSD-APDB Evaluates Multiple Sequence Alignment using structural information with APDB and iRMSD. >> Cite

T-RMSD Allows fine-grained structural clustering of a given group of related protein domains WEWM >> Cite
Other

Advanced Run your alignment using full featured T-Coffee options. >> Cite

## Constrained MSA

| H－RNase3 |  |
| :---: | :---: |
| H－RNase2 | －KPPEF＇TWA ONFE T HHINMTSQ d d inamQ |
| BP－RNaseA |  |
| BS－RNase |  |
| H－RNaseA |  |
| H－RNase 4 | －－MQDGMY QRFLR Of VHPEET－GGSDR 2 ¢NLMMQR |
| RC－RNase | QNWATEQQRAIINTPIIN－－－－UNTIMDN |
| H－RNase 3 |  |
| H－RNase2 |  |
| BP－RNaseA |  |
| BS－RNase |  |
| H－RNaseA |  |
| H－RNase 4 |  |
| RC－RNase | NIYIVGGgdakVNTFIISSATTVKAICTG－－VINMN－－－－－－VLSTTREQLNIGIRTSI |
| H－RNase 3 |  |
| H－RNase2 | SPQNISNCRYAQTPANMFYIVP他NRDQRRDPPQYPVVPY／ |
| BP－RNaseA | S－－KYPNC AYKTTQANKHIIVA¢EGN－－－－－－－－－PYVPYHEDASV－－－－ |
| BS－RNase | S－－KYPNCAYKTTQVEKHIIVAdGGK－－－－－－－－－PSVPY Fipasv－－－－ |
| H－RNaseA | S－－RYPNC AYRTSPKERHIIVA电GS－－－－－－－－－PYVPYHFDASVEDST |
| H－RNase 4 |  |
| RC－RNase | T－－－PRECYSSRTETNYICVKGEN－－－－－－－－－－－－YPV年AGIGRCP－ |

Fig．1．The multiple sequence alignment of seven RNases by WorkBench 3．2：The key active site residues homologous to His12，Lys41，and His119 of BP－RNaseA，the cysteine residues responsible for disulfide bond linkage and two matched Gln residues are shown in boxes．


Fig．2．The multiple sequence alignment of seven RNases by our CMSA：The key active site residues homologous to His12，Lys41，and His119 of BP－RNaseA，the cysteine residues responsible for disulfide bond linkage，and two matched Gln residues are shown in boxes．

## Homology extension approach

Chang, J.-M., Tommaso, P., Taly, J.-F. \& Notredame, C. Accurate multiple sequence alignment of transmembrane proteins with PSI-Coffee. Bmc Bioinformatics 13, 1-7 (2012).

## Homology-extended

Que1: how to build a profile?

## Que2: how to score profiles?



## Searching parameters

- Fast, Insensitive search
- High percent identity
- blastp -F "m S" -f 999 -M BLOSUM80 -G 9 -E 2 -e 1e-5
- Slow, Sensitive search
- Increase sensitivity, decrease specificity
- blastp -F "m S" -f 9 -M BLOSUM45 -e 100 -b 10000 -v 10000
- BLAST, page 146, 147
- By Ian Korf, Joseph Bedell, Mark Yandell
- Publisher: O'Reilly Media
- Release Date: July 2003



## Database Size


keyword:"Transmembrane [KW-0812]"

| Data Set | No. |
| :--- | ---: |
| UniRef50-TM | 87,989 |
| UniRef90-TM | 263,306 |
| UniRef100-TM | 613,015 |
| UniProt-TM | 818,635 |
| UniRef50 | $3,077,464$ |
| UniRef90 | $6,544,144$ |
| UniRef100 | $9,865,668$ |
| UniProt | $11,009,767$ |
| NCBI NR | $10,565,004$ |

## Performance comparison of different database sizes for the BAliBASE2-ref7.

- UniRef50-TM contains about 100 times fewer sequences than the full UniProt.
- The level accuracy is comparable and even superior to that achieved with the default PSI-Coffee while the CPU time requirements are dramatically decreased by a factor 10 .

| database | \# of seqs | SP | TC | extension(s) | total(s) |
| :--- | ---: | ---: | ---: | ---: | ---: |
| default T-Coffee | 0 | 0.911 | 0.498 | 0 | 2,735 |
| UniRef50-TM | 87,989 | 0.916 | 0.561 | 1,483 | 8,177 |
| UniRef90-TM | 263,306 | 0.918 | 0.548 | 3,343 | 9,610 |
| UniRef100-TM | 613,015 | 0.925 | 0.545 | 6,499 | 12,111 |
| UniProt-TM | 818,635 | 0.923 | 0.536 | 7,871 | 13,285 |
| UniRef50 | $3,077,464$ | 0.920 | 0.553 | 19,087 | 26,442 |
| UniRef90 | $6,544,144$ | 0.924 | 0.561 | 40,448 | 46,478 |
| UniRef100 | $9,865,668$ | 0.922 | 0.554 | 66,696 | 71,895 |
| UniProt | $11,009,767$ | 0.923 | 0.563 | 66,964 | 72,199 |
| NCBI NR | $10,565,004$ | 0.921 | 0.554 | 65,201 | 70,375 |

## 2.Phylogenetic trees

2.1 Enumerating trees and selecting search strategies

## \# of rooted and unrooted trees: 3 OTUs



For three operational taxonomic units (OTUs) there is one possible unrooted tree.


Any of the three edges can be selected to form a root.


Three rooted trees are possible.

## \# of rooted and unrooted trees: 4 OTUs


(b)

For 4 OTUs there are three possible unrooted trees.

For 4 OTUs there are 15 possible rooted trees.

There is only one of these 15 trees that accurately describes the evolutionary process by which these four sequences evolved.

## $\operatorname{TU}(k)$ ：the \＃of unrooted tree for $n$ taxa



Let $E(k)$ denote the \＃of edges in the unrooted tree for $k$ species．

$$
E(k)=2 k-3
$$

Let $T U(k)$ denote the \＃of unrooted trees for $k$ species．

$$
\begin{aligned}
& T U(k)=T U(k-1) \times E(k-1) \\
& =(T U(k-2) \times E(k-2)) \times E(k-1) \\
& =\prod_{i=1}^{k-2} E(k-i) \\
& =\prod_{i=1}^{k-2}(2 k-2 i-3) \\
& =(2 k-5) \times \cdots \times 5 \times 3 \times 1
\end{aligned}
$$

## TU(k) function

$$
\begin{aligned}
& T U(k)=1 \times 3 \times 5 \times \cdots \times(2 k-5) \\
& =\frac{1 \times 2 \times 3 \times 4 \times 5 \times \cdots \times(2 k-6) \times(2 k-5)}{2 \times 4 \times \cdots \times(2 k-6)} \\
& =\frac{(2 k-5)!}{(2 \times 1) \times(2 \times 2) \cdots(2 \times(k-3))} \\
& =\frac{(2 k-5)!}{(2 \times 2 \times \cdots \times 2) \times(1 \times 2 \times \cdots \times(k-3))} \\
& =\frac{(2 k-5)!}{2^{k-3}(k-3)!}
\end{aligned}
$$

| $k$ | $T U(k)$ |
| :--- | ---: |
| 4 | 3 |
| 5 | 15 |
| 6 | 105 |
| 7 | 945 |
| 8 | 10395 |
| 9 | 135,135 |
| 10 | $2,027,025$ |
| $\ldots$ | $\ldots$ |
| 20 | $\sim 2 \times 10^{20}$ |

## $\operatorname{TR}(k)$ ：the \＃of rooted tree for $k$ species

$$
\begin{aligned}
& T R(k)=T U(k) \times(2 k-3) \\
& =T U(k) \times E(k) \\
& =T U(k+1)
\end{aligned}
$$

| $k$ | $T U(k)$ | $\operatorname{TR}(k)$ |
| ---: | ---: | ---: |
| 2 | 1 | 1 |
| 3 | 1 | 3 |
| 4 | 3 | 15 |
| 5 | 15 | 105 |
| 10 | $2,027,025$ | $34,459,425$ |
| 20 | $\sim 2 \times 10^{20}$ | $\sim 8 \times 10^{21}$ |

## Stage 1: Use of DNA, RNA, or protein

- For phylogeny, DNA can be more informative.
- Some substitutions in a DNA sequence alignment can be directly observed: single nucleotide substitutions, sequential substitutions, coincidental substitutions.
- Additional mutational events can be inferred by analysis of ancestral sequences.

Two sequences (human and mouse) and their common ancestor: we can infer which DNA changes occurred over time


| ancestral | M | V | H | L | S | P | V | E | K | S | A | V |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| human | M | V | H | L | T | P | E | E | K | S | A | V |
| mouse | M | V | H | L | T | D | A | E | K | S | A | V |

## protein

DNA
human $5^{\prime}$ ATG GTG CAT CTG ACT CCT GAG GAG AAG TCT GCC GTT 3'
mouse 5. ATG GTG CAC CTG ACT GAT GCT GAG, AAG TCT GCT GTC 3.

| A | A | A | AA |  |
| :---: | :---: | :---: | :---: | :---: |
| G | $\mathrm{G} \rightarrow \mathrm{C}$ | $\mathrm{G} \rightarrow \mathrm{C}$ | CC | parallel substitutions |
| T | T | T | TT |  |
| c | c | $C \rightarrow G$ | CG | single substitution |
| C | C | $\mathrm{C} \rightarrow \mathrm{T} \rightarrow \mathrm{A}$ | CA | sequential substitution |
| T | T | T | TT |  |
| G | G | G | GG |  |
| T | $\mathrm{T} \rightarrow \mathrm{A}$ | $T \rightarrow \mathrm{C}$ | AC | coincidental substitutions |
| T | $T \rightarrow G$ | T | GT |  |
| C | $C \rightarrow G$ | $\mathrm{C} \rightarrow \mathrm{T} \rightarrow \mathrm{G}$ | GG | convergent substitutions |
| A | A | A | AA |  |
| G | $\mathrm{G} \rightarrow \mathrm{T} \rightarrow \mathrm{G}$ | G | GG | back substitution |

## Stage 2: Multiple sequence alignment

I. Confirm that all sequences are homologous
2. Adjust gap creation and extension penalties as needed to optimize the alignment
3. Restrict phylogenetic analysis to regions of the multiple sequence alignment for which data are available for all taxa (delete columns having incomplete data).

## Stage 3: models of DNA and amino acid substitution

- The simplest approach to measuring distances between sequences
- align pairs of sequences
- count the number of differences.
- For an alignment of length $N$ with $n$ sites at which there are differences, the degree of divergence $D$ is (Hamming distance)
- $D=n / N$
- But observed differences do not equal genetic distance!
- Genetic distance involves mutations that are not observed directly.


## Step matrices: number of steps required to change a character

(a)

|  | $A$ | $C$ | $T$ | $G$ |
| :---: | :---: | :---: | :---: | :---: |
| $A$ | 0 | 1 | 1 | 1 |
| $C$ | 1 | 0 | 1 | 1 |
| T | 1 | 1 | 0 | 1 |
| G | 1 | 1 | 1 | 0 |

(b)

|  | A | C | D | E | F | G | H | I | K | L | M | N | P | Q | R | S | T | V | W | Y |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | 0 | 2 | 1 | 1 | 2 | 1 | 2 | 2 | 2 | 2 | 2 | 2 | 1 | 2 | 2 | 1 | 1 | 1 | 2 | 2 |
| C |  | 0 | 2 | 3 | 1 | 1 | 2 | 2 | 3 | 2 | 3 | 2 | 2 | 3 | 1 | 1 | 2 | 2 | 1 | 1 |
| D |  |  | 0 | 1 | 2 | 1 | 1 | 2 | 2 | 2 | 3 | 1 | 2 | 2 | 2 | 2 | 2 | 1 | 3 | 1 |
| E |  |  |  | 0 | 3 | 1 | 2 | 2 | 1 | 2 | 2 | 2 | 2 | 1 | 2 | 2 | 2 | 1 | 2 | 2 |
| F |  |  |  |  | 0 | 2 | 2 | 1 | 3 | 1 | 2 | 2 | 2 | 3 | 2 | 1 | 2 | 1 | 2 | 1 |
| G |  |  |  |  |  | 0 | 2 | 2 | 2 | 2 | 2 | 2 | 2 | 2 | 1 | 1 | 2 | 1 | 1 | 2 |
| H |  |  |  |  |  |  | 0 | 2 | 2 | 1 | 3 | 1 | 1 | 1 | 1 | 2 | 2 | 2 | 3 | 1 |
| I |  |  |  |  |  |  |  | 0 | 1 | 1 | 1 | 1 | 2 | 2 | 1 | 1 | 1 | 1 | 3 | 2 |
| K |  |  |  |  |  |  |  |  | 0 | 2 | 1 | 1 | 2 | 1 | 1 | 2 | 1 | 2 | 2 | 2 |
| L |  |  |  |  |  |  |  |  |  | 0 | 1 | 2 | 1 | 1 | 1 | 1 | 2 | 1 | 1 | 2 |
| M |  |  |  |  |  |  |  |  |  |  | 0 | 2 | 2 | 2 | 1 | 2 | 1 | 1 | 2 | 3 |
| N |  |  |  |  |  |  |  |  |  |  |  | 0 | 2 | 2 | 2 | 1 | 1 | 2 | 3 | 1 |
| P |  |  |  |  |  |  |  |  |  |  |  |  | 0 | 1 | 1 | 1 | 1 | 2 | 2 | 2 |
| Q |  |  |  |  |  |  |  |  |  |  |  |  |  | 0 | 1 | 2 | 2 | 2 | 2 | 2 |
| R |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 0 | 1 | 1 | 2 | 1 | 2 |
| S |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 0 | 1 | 2 | 1 | 1 |
| T |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 0 | 2 | 2 | 2 |
| V |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 0 | 2 | 2 |
| W |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 0 | 2 |
| Y |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 0 |

nucleotide step matrix
amino acid step matrix For amino acids, between 1 and 3 nucleotide changes are required to change one residue to another.

## Quantification of evolutionary distances

## Evolutionary Distances

- They measure the total number of substitutions that occurred on both lineages since divergence from last common ancestor.
- Divided by sequence length.
- Expressed in substitutions / site
ancestor

sequence 1 sequence 2


## The problem of hidden or multiple changes

- $D$ (true evolutionary distance) $\geq$ fraction of observed differences ( $p$ )

- $D=p+$ hidden changes
- Through hypotheses about the nature of the residue substitution process, it becomes possible to estimate $D$ from observed differences between sequences.


## Correcting for multiple substitutions



## Correcting for multiple substitutions

- Requires a statistical 'model' of how the process of substitution works to correct for
- Differences in the rates of different substitution types
- Jukes and Cantor - all substitutions are treated the same
- Kimura 2-parameter model - distinguishes between transitions and transversions
- Different frequencies of different nucleotides
- GC content - the HKY model adds nucleotide frequency parameters to the Kimura 2-parameter model
- Different rates at different sites (often modelled using a distribution - e.g. Gamma distribution)


## Stage 3: Jukes and Cantor one-parameter model of nucleotide substitution

- This model describes the probability that one nucleotide will change into another. It assumes that each residue is equally likely to change into any other.
- Jukes and Cantor (I969) proposed a corrective formula:

$$
D=\left(-\frac{3}{4}\right) \ln \left(l-\frac{4}{3} p\right)
$$



## JC model: $D=\left(-\frac{3}{4}\right) \ln \left(1-\frac{4}{3} p\right)$

- Consider an alignment where $3 / 60$ aligned residues differ
- The normalized Hamming distance, $3 / 60=0.05$.
- The Jukes-Cantor correction is

$$
D=\left(-\frac{3}{4}\right) \ln \left(1-\frac{4}{3} 0.05\right)=0.052
$$

- When 30/60 aligned residues differ, the Jukes-Cantor correction is more substantial:

$$
D=\left(-\frac{3}{4}\right) \ln \left(1-\frac{4}{3} 0.5\right)=0.82
$$

## Two DNA substitution mutations

- Transitions: interchanges of two-ring purines or of one-ring pyrimidines: they therefore involve bases of similar shape.
- A $<->$ G, C $<->$ T
- Transversions: interchanges of purine for pyrimidine bases, which therefore involve exchange of nne-rino and twn-ring structures.
- $\mathrm{A} \leftrightarrow->\mathrm{C}, \mathrm{A} \leftrightarrow->\mathrm{T}, \mathrm{G}\langle-\mathrm{C}, \mathrm{G}\langle->\mathrm{T}$


Kimura two-parameter model of nucleotide substitution (assumes a $\neq b$ )


## Kimura's two parameter distance (DNA)

- Hypotheses of the model
- All sites evolve independently and following the same process.
- Substitutions occur according to two probabilities
- Transitions : G <->A or C <->T
- Transversions : other changes
- The base substitution process is constant in time.
- Quantification of evolutionary distance (d) as a function of the fraction of observed differences $p$ : transitions, $q$ : transversions
- $d=-\frac{1}{2} \ln [(1-2 p-q) \sqrt{1-2 q}]$


## There are dozens of models

Jukes-Cantor model
(a)


Kimura model


Tamura model


## B\&FG 3e

## Substitution model categories



# Stage 4: tree-building methods 

distance-based
maximum parsimony
maximum likelihood
Bayesian methods

## Main families of Methods for Phylogenetic reconstruction

## COMPUTATIONAL METHOD



## UPGMA

Unweighted Pair-Group Method with Arithmetic Mean

## Tree-building methods: UPGMA

- Step 2: Find the two proteins with the smallest pairwise distance. Cluster them.



## Tree-building methods: UPGMA

- Step 3: Do it again. Find the next two proteins with the smallest pairwise distance. Cluster them.



## Tree-building methods: UPGMA

- Step 4: Keep going. Cluster.



## Tree-building methods: UPGMA

- Step 4: Last cluster! This is your tree.

(a)


|  | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | - |  |  |  |  |
| 2 | 0.1 | - |  |  |  |
| 3 | 0.8 | 0.8 | - |  |  |
| 4 | 0.8 | 1 | 0.3 | - |  |
| 5 | 0.9 | 0.9 | 0.3 | 0.2 | - |

(b)

(c)

(d)

(e)


## Distance-based methods: UPGMA trees

- UPGMA is a simple approach for making trees.
- An UPGMA tree is always rooted.
- An assumption of the algorithm is that the molecular clock is constant for sequences in the tree. If there are unequal substitution rates, the tree may be wrong.
- While UPGMA is simple, it is less accurate than the neighbor-joining approach (described next).


## Neighbor-Joining

N Saitou, M Nei, The neighbor-joining method: a new method for reconstructing phylogenetic trees. Molecular Biology and Evolution, Volume 4, Issue 4, Jul 1987, Pages 406-425

## Saitou, N. \& Nei, M. Mol. Biol. Evol. 4, 406-425 (1987).

Number 20 on the list is a paper ${ }^{12}$ that introduced the "neighbor-joining" method, a fast, efficient way of placing a large number of organisms into a phylogenetic tree according to some measure of evolutionary distance between them, such as genetic variation.

It links related organisms together one pair at a time until a tree is resolved. Physical anthropologist Naruya Saitou helped to devise the technique when he joined Masatoshi Nei's lab at the University of Texas in Houston in the 1980s to work on human evolution and molecular genetics, two fields that were starting to burst at the seams with information.

Saitou, N. \& Nei, M. Mol. Biol. Evol. 4, 406-425 (1987)

Another field buoyed by the growth in genome sequencing is phylogenetics, the study of evolutionary relationships between species.
"We physical anthropologists were facing kind of the big data of that time," says Saitou, now at Japan's National Institute of Genetics in Mishima. The technique made it possible to devise trees from large data sets without eating up computer resources. (And, in a nice cross-fertilization within the top-10, Clustal's algorithms use the same strategy.)

## Why NJ instead of UPGMA?

In the original CLUSTAL programs, the initial guide trees, used to guide the multiple alignment, were calculated using the UPGMA method.

We now use the Neighbour-Joining method which is more robust against the effects of unequal evolutionary rates in different lineages and which gives better estimates of individual branch lengths.

This is useful because it is these branch lengths which are used to derive the sequence weights.

We also allow users to choose between fast approximate alignments or full dynamic programming for the distance calculations used to make the guide tree.

## Making trees using neighbor-joining

- useful for making a tree having a large number of taxa.
- Begin by placing all the taxa in a star-like structure.



## The algorithm

- Based on the current distance matrix calculate the matrix $Q$.
- Find the pair of distinct taxa i and j for which $Q(i, j)$ has its lowest value. These taxa are joined to a newly created node, which is connected to the central node.
- Calculate the distance from each of the taxa in the pair to this new node.
- Calculate the distance from each of the taxa outside of this pair to the new node.
- Start the algorithm again, replacing the pair of joined neighbors with the new node and using the distances calculated in the previous step.


## The algorithm

1. Based on the current distance matrix calculate the matrix $Q$.
2. Find the pair of distinct taxa $i$ and $j$ for which $Q(i, j)$ has its lowest value. These taxa are joined to a newly created node, which is connected to the central node.
3. Calculate the distance from each of the taxa in the pair to this new node.
4. Calculate the distance from each of the taxa outside of this pair to the new node.
5. Start the algorithm again, replacing the pair of joined neighbors with the new node and using the distances calculated in the previous step.

## The matrix $Q$

1. Based on the current distance matrix $D$ calculate the matrix $Q$.

- $Q(i, j)=(n-2) d(i, j)-\sum_{k=1}^{n} d(i, k)-\sum_{k=1}^{n} d(n, k)$

| D | $\mathbf{a}$ | $\mathbf{b}$ | $\mathbf{c}$ | $\mathbf{d}$ | $\mathbf{e}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{a}$ | 0 | 5 | 9 | 9 | 8 |
| $\mathbf{b}$ | 5 | 0 | 10 | 10 | 9 |
| $\mathbf{c}$ | 9 | 10 | 0 | 8 | 7 |
| $\mathbf{d}$ | 9 | 10 | 8 | 0 | 3 |
| $\mathbf{e}$ | 8 | 9 | 7 | 3 | 0 |


| $\mathbf{Q}_{1}$ | $\mathbf{a}$ | $\mathbf{b}$ | $\mathbf{c}$ | $\mathbf{d}$ | $\mathbf{e}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{a}$ |  | -50 | -38 | -34 | -34 |
| $\mathbf{b}$ | -50 |  | -38 | -34 | -34 |
| $\mathbf{c}$ | -38 | -38 |  | -40 | -40 |
| $\mathbf{d}$ | -34 | -34 | -40 |  | -48 |
| $\mathbf{e}$ | -34 | -34 | -40 | -48 |  |

Credit by https://en.wikipedia.org/wiki/Neighbor joining

## Join two nodes

2. Find the pair of distinct taxa $i$ and $j$ for which $Q(i, j)$ has its lowest value. These taxa are joined to a newly created node, which is connected to the central node.

- Merge nodes $a$ and $b$ into $u$

| $Q_{1}$ | $\mathbf{a}$ | $\mathbf{b}$ | $\mathbf{c}$ | $\mathbf{c}$ | $\mathbf{e}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{a}$ |  | -50 | -38 | -34 | -34 |
| $\mathbf{b}$ | -50 |  | -38 | -34 | -34 |
| $\mathbf{c}$ | -38 | -38 |  | -40 | -40 |
| $\mathbf{d}$ | -34 | -34 | -40 |  | -48 |
| $\mathbf{e}$ | -34 | -34 | -40 | -48 |  |



Credit by https://en.wikipedia.org/wiki/Neighbor joining

## Distance from the new node

- 3. Calculate the distance from each of the taxa in the pair to this new node. i.e, Merge nodes $f$ and $g$ into $u$
- $\delta(f, u)=\frac{1}{2} d(f, g)+\frac{1}{2(n-2)}\left[\sum_{k=1}^{n} d(f, k)-\sum_{k=1}^{n} d(g, k)\right]$
- $\delta(g, u)=d(f, g)-\delta(f, u)$
- Example: Merge nodes $a$ and $b$ into $u$
- $\delta(a, u)=\frac{1}{2} d(a, b)+\frac{1}{2(5-2)}\left[\sum_{k=1}^{5} d(a, k)-\sum_{k=1}^{5} d(b, k)\right]=\frac{5}{2}+\frac{31-34}{6}=2$
- $\delta(b, u)=d(a, b)-\delta(a, u)=5-2=3$

|  | $\mathbf{a}$ | $\mathbf{b}$ | $\mathbf{c}$ | $\mathbf{d}$ | $\mathbf{e}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{a}$ | 0 | 5 | 9 | 9 | 8 |
| $\mathbf{b}$ | 5 | 0 | 10 | 10 | 9 |
| $\mathbf{c}$ | 9 | 10 | 0 | 8 | 7 |
| $\mathbf{d}$ | 9 | 10 | 8 | 0 | 3 |
| $\mathbf{e}$ | 8 | 9 | 7 | 3 | 0 |



## Distance of the other taxa from the new node

- 4. Calculate the distance from each of the taxa outside of this pair to the new node.
- $d(u, k)=\frac{1}{2}[d(f, k)+d(g, k)-d(f, g)]$
- Example: Merge nodes $a$ and $b$ into $u$
- $d(u, c)=\frac{1}{2}[d(a, c)+d(b, c)-d(a, b)]=\frac{9+10-5}{2}=7$
- $d(u, d)=\frac{1}{2}[d(a, d)+d(b, d)-d(a, b)]=\frac{9+10-5}{2}=7$
- $d(u, e)=\frac{1}{2}[d(a, e)+d(b, e)-d(a, b)]=\frac{8+9-5}{2}=6$

|  | $\mathbf{a}$ | $\mathbf{b}$ | $\mathbf{c}$ | $\mathbf{d}$ | $\mathbf{e}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{a}$ | 0 | 5 | 9 | 9 | 8 |
| $\mathbf{b}$ | 5 | 0 | 10 | 10 | 9 |
| $\mathbf{c}$ | 9 | 10 | 0 | 8 | 7 |
| $\mathbf{d}$ | 9 | 10 | 8 | 0 | 3 |
| $\mathbf{e}$ | 8 | 9 | 7 | 3 | 0 |


|  | $u$ | $c$ | $d$ | $e$ |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{u}$ | 0 | 7 | 7 | 6 |
| $c$ | 7 | 0 | 8 | 7 |
| $d$ | 7 | 8 | 0 | 3 |
| $e$ | 6 | 7 | 3 | 0 |

## Repeat

5. Start the algorithm again, replacing the pair of joined neighbors with the new node and using the distances calculated in the previous step.

|  | $\mathbf{u}$ | c | d | e |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{u}$ | 0 | 7 | 7 | 6 |
| $\mathbf{c}$ | 7 | 0 | 8 | 7 |
| $\mathbf{d}$ | 7 | 8 | 0 | 3 |
| $\mathbf{e}$ | 6 | 7 | 3 | 0 |


| $\mathrm{Q}_{2}$ | $\mathbf{u}$ | $\mathbf{c}$ | $\mathbf{d}$ | $\mathbf{e}$ |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{u}$ |  | -28 | -24 | -24 |
| $\mathbf{c}$ | -28 |  | -24 | -24 |
| $\mathbf{d}$ | -24 | -24 |  | -28 |
| $\mathbf{e}$ | -24 | -24 | -28 |  |



## Maximum Parsimony Method

## Tree-building methods: character based

- Rather than pairwise distances between proteins, evaluate the aligned columns of amino acid residues (characters).
- Tree-building methods based on characters include
- maximum parsimony
- maximum likelihood


## Maximum Parsimony (MP)

- Find the tree with the shortest branch lengths possible. Thus we seek the most parsimonious ("simple") tree.
- Identify informative sites.
- Constant characters are not parsimony-informative.
- Construct trees, counting the number of changes required to create each tree.
- <= 12 taxa : evaluate all possible trees exhaustively
- >12 taxa : perform a heuristic search.
- Select the shortest tree (or trees).


## An example of tree-building using MP

- Consider these four taxa

AAG
AAA
GGA
AGA

- How might they have evolved from a common ancestor such as AAA?


## MP

- Choose the tree(s) with the lowest cost (shortest branch lengths)


Cost $=3$
Cost $=4$
Cost $=4$

## Select the tree supported by the largest number of Informative Site



# Maximum Likelihood Method 

## Making trees using maximum likelihood

- An alternative to maximum parsimony.
- What are the tree topology and branch lengths that have the greatest likelihood of producing the observed data set?
- ML is implemented in the TREE-PUZZLE program, as well as MEGA5, PAUP and PHYLIP.


## Likelihood

- Given some data ( $D$ ) a decision must be made about an adequate explanation (H, hypothesis)
- D: alignment
- H: Model of evolution, tree topology, branch lengths, parameters of the model
- $L=\operatorname{Pr}(D \mid H)$
- Each $H$ will have a certain probability of producing the data


## Likelihood vs Probability

- https://youtu.be/pYxNSUDSFH4
- The likelihood function != the probability of a hypothesis being correct!
- The likelihood function is defined in terms of probability of producing the observed events not of the unknown parameters
- Thus: the probability of observing the data has nothing to do with the probability that the underlying model is correct.


## Maximum Likelihood

- https://youtu.be/XepXtI9YKwc
- Given some data ( $D$ ) a decision must be made about an adequate explanation ( $H$, hypothesis)
- $L=\operatorname{Pr}(D \mid H)$
- Each $H$ will have a certain probability of producing the data
- The best $H$ is that of the greatest $P$


## Coin Example

- Data: flipping coins and counting the number of times "heads" appear
- You throw the coin twice and observe "heads" both times.
- Hypotheses : You might have two hypotheses to explain these data.
- $H_{1}$, the coin is normal: $p=0.5$, of appearing head.
- $\mathrm{H}_{2}$, the coin is rigged with an $80 \%$ chance of getting a head, $p$ $=0.8$.
- What is the likelihood of $\mathrm{H}_{1}$ ?
- What is the likelihood of $\mathrm{H}_{2}$ ?


## Likelihood of the coin example

- The probability of observing "heads" in each of two flips
- under $\mathrm{H} 1, \mathrm{~L}($ data $\mid \mathrm{H} 1)=0.5 \times 0.5=0.25$
- under H2, L(data|H2) $=0.8 \times 0.8=0.64$
- Since the probability of observing the data under H 2 is greater than under H 1 , you might argue that the "rigged" coin hypothesis is the more likely.


## Parameter Estimation

- Assuming sample $x_{1}, x_{2}, \ldots, x_{n}$ is from a parametric distribution $f(x \mid \theta)$, estimate $\theta$.
- Given sample HHTHH of (possibly biased) coin flips, estimate $\theta=$ probability of Heads
- $\operatorname{Pr}(\mathrm{HHTHH} \mid .6)>\operatorname{Pr}(\mathrm{HHTHH} \mid .5)$, event HHTHH is more likely when $\theta=.6$ than $\theta=.5$
- And what $\theta$ make HHTHH most likely?


## Likelihood Function

- Probability of HHTHH, given $\theta$ :

| $\boldsymbol{\theta}$ | $\boldsymbol{0}^{4}(1-\boldsymbol{\theta})$ |
| :---: | :---: |
| 0.20 | 0.0013 |
| 0.50 | 0.0313 |
| 0.80 | 0.0819 |
| 0.95 | 0.0407 |



## Maximum Likelihood Parameter Estimation

- As a function of $\theta$, what $\theta$ maximizes the likelihood of the data actually observed by taking derivative of $L$ (Pr) with respect to $\theta$

$$
\frac{d L}{d \theta}=4 \theta^{3}-5 \theta^{4}=\theta^{3}(4-5 \theta)
$$

- equating to 0 , and solving

$$
\frac{d L}{d \theta}=4 \theta^{3}-5 \theta^{4}=\theta^{3}(4-5 \theta)=0 \rightarrow \theta=\frac{4}{5}
$$

- More easily, likelihood are often maximized by maximizing their logarithm

$$
\ln L=4 \ln \theta+\ln (1-\theta)
$$

- whose derivative is

$$
\frac{d \ln L}{d \theta}=\frac{4}{\theta}-\frac{1}{1-\theta}=0 \rightarrow \theta=\frac{4}{5}
$$

## First use in phylogenetics

- Cavalli-Sforza and Edwards (1967) for gene frequency data
- Felsenstein (1981) for DNA sequences
- In phylogenetics, the hypothesis is
- a tree topology
- its branch-lengths
- a model under which the data evolved



## Maximum Likelihood Method(con’t)

- $s$ homologous sequences each with $N$ nucleotides
- $X_{k}=\left(\mathrm{X}_{1 \mathrm{k}}, \ldots, X_{\text {sk }}\right)$ the nucleotide configuration at kth site
- The likelihood function of tree $T$ at the $k$ th site
- The likelihood function for the entire sequence for tree $T$

$$
L\left(\theta_{1}, \ldots, \theta_{\eta} \mid X_{1}, \ldots, X_{N}, T\right)=\prod_{k=1}^{N} f\left(X_{k} \mid \theta, T\right)
$$

## An example

- The model is reversible, ie. $p(A \rightarrow G)=p(A \rightarrow G)$, so the root can be placed at any node
- Pattern probability $=\mathrm{p}(\mathrm{G} \rightarrow \mathrm{G}) \times \mathrm{p}(\mathrm{G} \rightarrow \mathrm{G}) \times \mathrm{p}(\mathrm{G} \rightarrow \mathrm{A})$ $\times p(A \rightarrow A) \times p(A \rightarrow A)$



## Site pattern probability

- Under the simple Jukes-Cantor model, all base frequencies=0.25, all substitutions equally probable.
- $P_{i j}=\left\{\begin{array}{l}0.25+0.75 e^{-b}, i=j \\ 0.25-0.75 e^{-b}, i \neq j\end{array}\right.$, where $b$ is branch-length
(subs/site)
- Where $b=0.5, \mathrm{P}_{\mathrm{ij}}(\mathrm{i}=\mathrm{j})=0.7049, \mathrm{P}_{\mathrm{ij}}(\mathrm{i} \neq \mathrm{j})=0.0984$
- Site pattern probability

$$
\begin{aligned}
& =p(G \rightarrow G) \times p(G \rightarrow G) \times p(G \rightarrow A) \times p(A \rightarrow A) \times p(A \rightarrow A) \\
& =0.7049 \times 0.7049 \times 0.0984 \times 0.7049 \times 0.7049 \\
& =0.0243
\end{aligned}
$$

A


The likelihood of a tree

- The likelihood of a tree = the product of the site likelihoods
- Taken as natural logs, the site likelihoods can be summed to give the log likelihood
- The sum of the probabilities for the 16 possible site patterns $=0.0333$

- Hence, the site $-\ln L=3.402$

the ML tree with the highest likelihood
- The tree with the highest likelihood (lowest $-\ln \angle$ )
- Tree 2 is the ML tree by $8.801-\ln L$ units( $=2052.456-2043.655)$


| Site | $-\ln L(1)$ | $-\ln L(2)$ |
| :--- | :---: | :---: |
| 1 | 2.457 | 2.891 |
| 2 | 1.568 | 1.943 |
| . | $\cdots$ | $\cdots$ |
| . | $\cdots$ | . |
| 1206 | 2.541 | 1.943 |
|  | 2052.456 | 2043.655 |

## Phylogenetic Relationship of CoVs



Zhang, T., Wu, Q. \& Zhang, Z. Probable Pangolin Origin of SARS-CoV-2 Associated with the COVID-19 Outbreak. Curr Biology Cb 30, 1346-1351.e2 (2020).

# If you build a tree, 



Clustal


Dalign


DCA


Mafft


Muscle


ProbCons


T-Coffee

Which guy should I trust?
panel D). Most studies ignore that these scores are based on a fixed sequence alignment that supports the tree in the first place; they may thus make us overly confident of its accuracy.

Ari Löytynoja and Nick Goldman, "Uniting Alignments and Trees," Science 324, no. 5934 (June 19, 2009): 1528-1529.

## YPL077C with six topologies



MAFFT (0.18)
they have lower bootstraps
PROBCONS (0.05)


Fig. 1. An example, involving ORF YPL077C, in which alignments produced by seven different alignment methods produce six different estimated trees, albeit with low bootstrap support (bootstrap proportions shown parenthetically for each tree).

Rank: 41 Citations: 21,373
Confidence limits on phylogenies: an approach using the bootstrap
Felsenstein, J.
Evolution 39, 783-791 (1985).



## Super multiple sequence alignment (SMSA)

Clustal

MAFFT
T-Coffee
S. cer
S. par
S. klu
S. cer
S. par
S. klu


## Andrey Zharkikh and Wen-Hsiung Li

- for Demographic and Population Genetics, University of P.O. Box 20334, Houston, Texas 77225
partial


Average bootstrap, AUC values and the number of TPs for 10 and 25 accepted FPs of each method

| Method | ave. Bootstrap | AUC | TPs for 10 FPs | TPs for 25 FPs | TPs in total |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Clustal | 51.31 | 0.7521 | 185 | 274 | 643 |
| DCA | 50.62 | 0.7694 | 194 | 284 | 624 |
| DIALIGN | 51.94 | 0.7618 | 253 | 340 | 659 |
| MAFFT | 52.82 | 0.7750 | 253 | 359 | 665 |
| Muscle | 52.35 | 0.7771 | 224 | 315 | 639 |
| Probnt | 50.96 | 0.7790 | 256 | 312 | 642 |
| T-Coffee | 51.21 | 0.7889 | 234 | 311 | 620 |
| M-Coffee | 51.41 | 0.7688 | 193 | 325 | 646 |
| SMSA | 77.31 | 0.8301 | 329 | 425 | 661 |
| pSMSA | 50.96 | 0.8140 | 342 | 385 | 661 |
| wpSMSA | 50.86 | 0.8215 | 353 | 423 | 661 |



## alignment uncertainty - data

OPOSSUM
BLOSUM62


## alignment uncertainty - data


alignment uncertainty - data

- It gets worse with a multiple sequence alignment.


Telling apart Uncertainty parts of the alignment is more important than the overall accuracy.

## Guidance



Penn O, Privman E, Landan G, Graur D, Pupko T (2010) An alignment confidence score capturing robustness to guide tree uncertainty. Mol Biol Evol 27: 1759-1767.

## Gblocks

a)

b)

d)


Talavera G, Castresana J (2007) Improvement of Phylogenies after

Sequence Alignments. Syst Biol 56: 564-577. ${ }^{\text {tation by }}$

## trimAl



Capella-Gutiérrez S, Silla-MartíRez ty Gabaldón T (2009) trimAl: a tool for automated alignment trimming intratgerprale phylogenetic analyses. Bioinformatics 25: 1972-1973.

consistency
inconsistency
inconsistency


## CLUSTAL W (1.83) multiple sequence alignment

1j46_A
2lef_A 1k99_A 1aab_

MQ------DRVKRP---MNAFIVWSRDQRRKMALENPRMRN--SEISKQL MH--------IKKP---LNAFMLYMKEMRANVVAESTLKES--AAINQIL MKKLKKHPDFPKKP---LTPYFRFFMEKRAKYAKLHPEMSN--LDLTKIL GK------GDPKKPRGKMSSYAFFVQTSREEHKKKHPDASVNFSEFSKKC

## Residue level

| Col | row | row | TCS |
| :--- | :--- | :--- | :--- |
| 1 | 1 | 2 | 0.762 |
| 1 | 1 | 3 | 0.748 |
| 1 | 1 | 4 | 0.741 |
| 1 | 2 | 3 | 0.651 |
| 1 | 2 | 4 | 0.677 |
| 1 | 3 | 4 | 0.693 |
| 2 | 1 | 3 | 0.562 |
| 2 | 1 | 4 | 0.632 |
| 2 | 3 | 4 | 0.526 |

T-COFFEE, Version_9.01 (2012-01-27 09:40:38) Cedric Notredame CPU TIME:0 sec.
SCORE=76 Alignment level

| BAD AVG | GOOD |
| :--- | :--- |
| * |  |
| 1j46_A | $:$ |
| 2lef_A | 74 |
| 1k99_A |  |
| 1aab_ | 75 |
| cons | $:$ |

1j46_A 75------4566---677777777777777777776666--7789999
2lef_A 6--------566---677777777777777777777766--7789999
1k99_A 865454445667---777788887888888888877877--7789999
1aab_ 76------5665333566676666666666666666655336789999

## Test2 - structural modeling @ alignment level

reference alignment
Guidence/TCS


SP1 - SP2 ? confidence1 - confidence2

Guidance $=71.10 \%$


[^1]
## Renewing Felsenstein’s phylogenetic bootstrap in the era of big data

- transfer distance, (b,b*) : a branch b of the reference tree T and a branch b* of a bootstrap tree T* is equal to the number of taxa that must be transferred (or removed), in order to make both branches identical
- Felsenstein (FBP) and transfer (TBE) bootstrap supports on the same tree with 9,147 HIV-1M pol sequences



## Significantly different output when changing sequence input order

- S-o-P comparison vs average identity (Spearman correlation rs $=0.79$ ).
- a high MSA structural accuracy variability vs correlating with MSA identity (rs = -0.51)




## Regressive algorithm enables MSA of up to 1.4 million sequences on


c


Merged MSA


# Special Issue "Phylogenetic Methods in the Genomic Era: Challenges in Multiple Sequence Alignment and Phylogenetics for Genome-Scale Data" 

## Guest Editors

Dr. Cedric Notredame
Pompeu Fabra University, Barcelona, Spain

Dr. Jia-Ming Chang
National Chengchi University, Taipei, Taiwan

Dr. Minh Bui
University of Melbourne, Melbourne, Australia

Dr. Ding He
University of Copenhagen,
Copenhagen, Denmark
Deadline for submissions:
1 June 2020

## Short Information about the Special Issue:

Grand-scale genome sequencing projects dedicating a systematic approach to targeting wellrecognized taxonomic groups have started to appear.

To highlight the importance of this exciting moment for phylogenetic method development and evolutionary data inference in facing the big data era, this Special Issue welcomes contributions of methods and metrics addressing challenges from sequence alignment to tree reconstruction in phylogenomics.

## Author Benefits

## $\delta$ Open Access

© No Copyright Constraints Retain copyright of your work and free use of your article

## ( ) Rapid Publication

No Space Constraints or Color Charges

## EO Thorough peer-review

## Thank you Ahy Qusestion?




[^0]:    Letko, M., Marzi, A. \& Munster, V. Functional assessment of cell entry and receptor usage for SARS-CoV-2 and other lineage B
    betacoronaviruses. Nat Microbiol5, 562-569 (2020).

[^1]:    $\square$ Ref. - MAFFT
    $\square$ Ref. - MUSCLE

    - Ref. - ClustalW
    - MAFFT - MUSCLE
    $\square$ MAFFT - ClustalW
    MUSCLE - ClustalW

